

Design and Analysis of a Content-Oriented Internet

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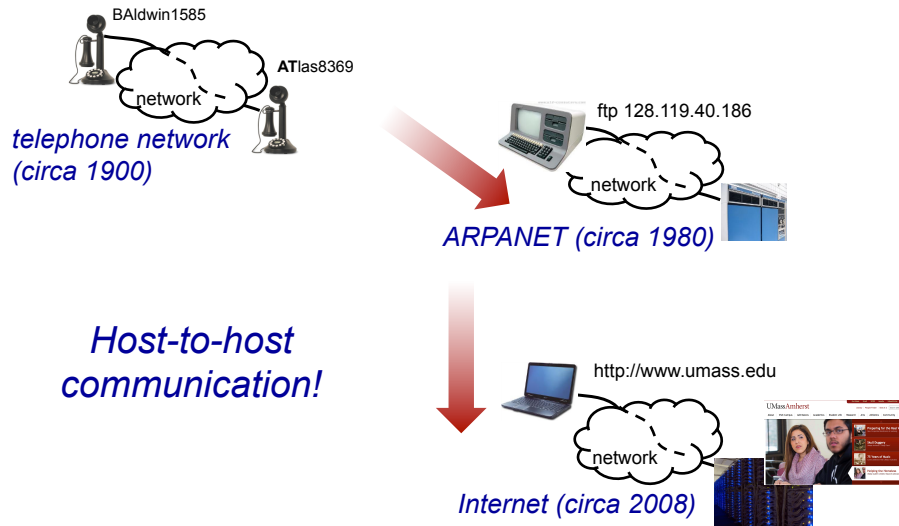
Overview

- ❖ introduction: “getting connected”
- ❖ Internet 101
- ❖ content-oriented networks:
 - “over the top” approaches
 - network-based approaches
- ❖ wrap-up

Goals:

- ❖ learn something about networking
- ❖ appreciate research challenges in content networking
 - learn about research here at UMass

Getting connected: a history



The Internet in 2014

- ❖ content is king!
 - video: 64% of all global consumer Internet traffic in 2012
 - TV, VOD, P2P: 80-90 % global consumer traffic by 2017
- ❖ *what not where*



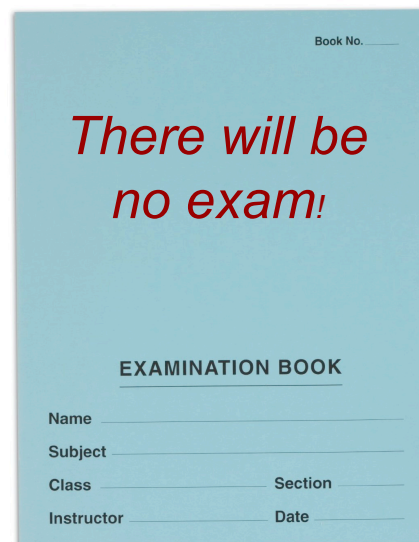
Challenge: can a network architecture designed for wired, host-host communication meet demands of today's (and tomorrow's) *content-oriented* mobile Internet?

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Internet 101

- ❖ hosts: clients, servers
- ❖ packets, routers
- ❖ Internet: a network of networks
- ❖ protocols



Internet I0I



billions of connected computing devices: *hosts*

- running *network apps*
- clients, servers

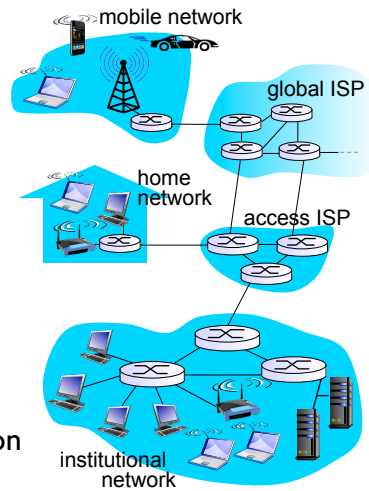


communication links

- wired, wireless
- transmission rate: bandwidth

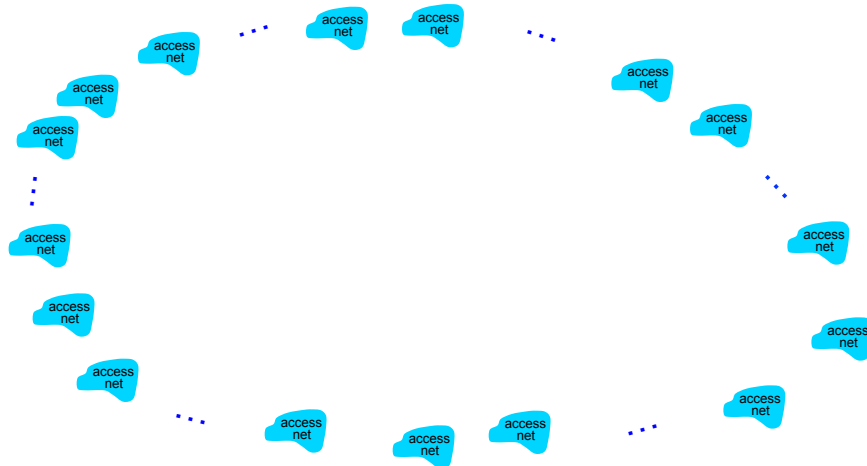


routers: switch/forward *packets* (chunks of data) from source to destination



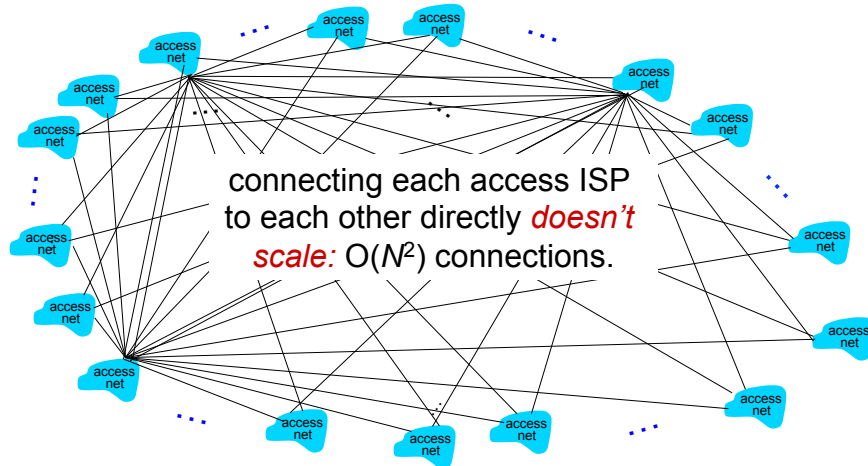
Internet I0I: a network of networks

Question: given *millions* of access ISPs, how to connect them together?



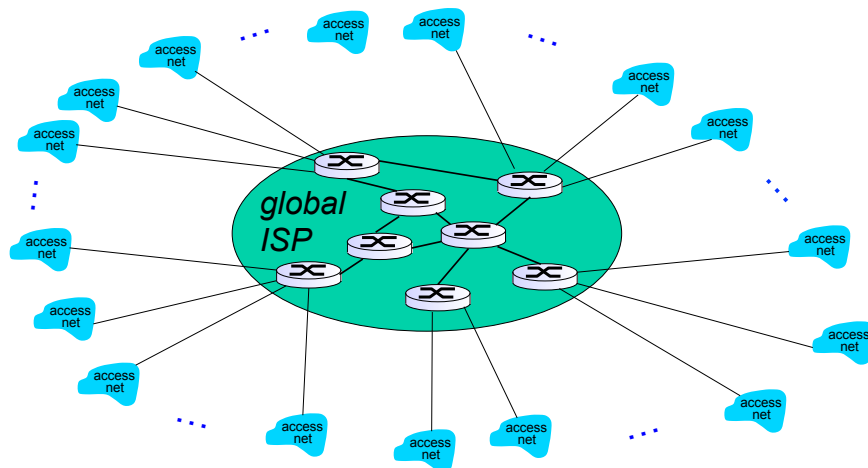
Internet I0I: a network of networks

Option: connect each access ISP to every other access ISP?



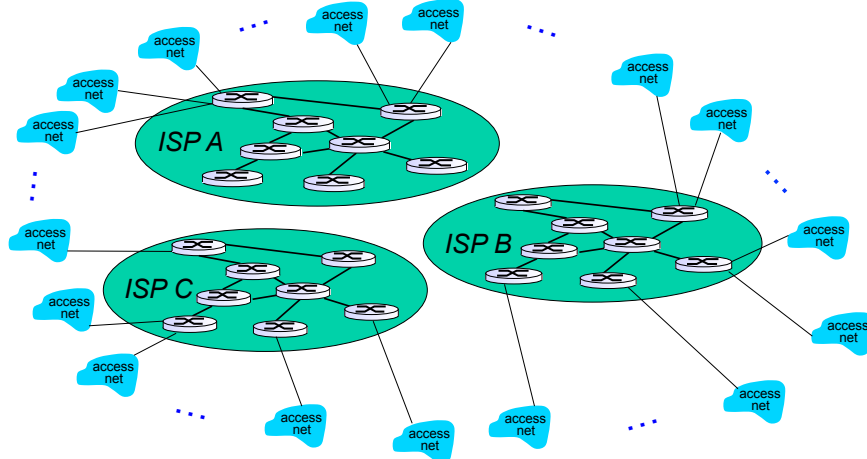
Internet I0I: a network of networks

Option: connect each access ISP to a global transit ISP?



Internet I0I: a network of networks

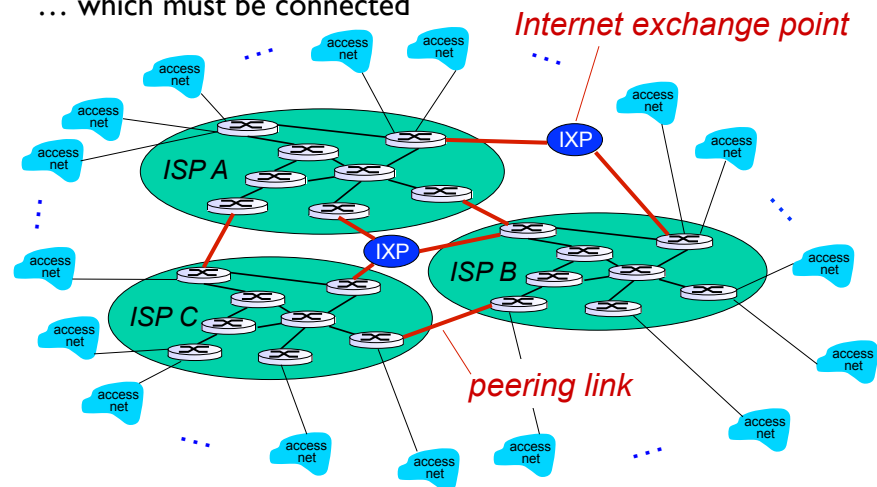
If one global ISP is viable business, there will be competitors



Internet I0I: a network of networks

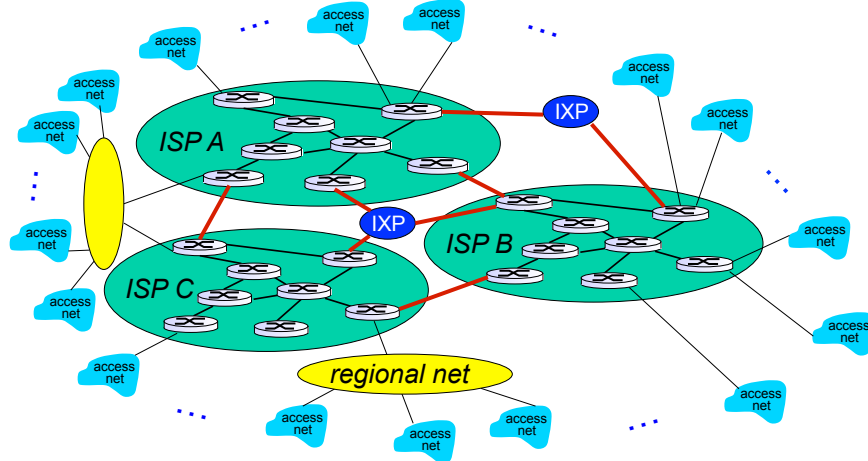
If one global ISP is viable business, there will be competitors

... which must be connected



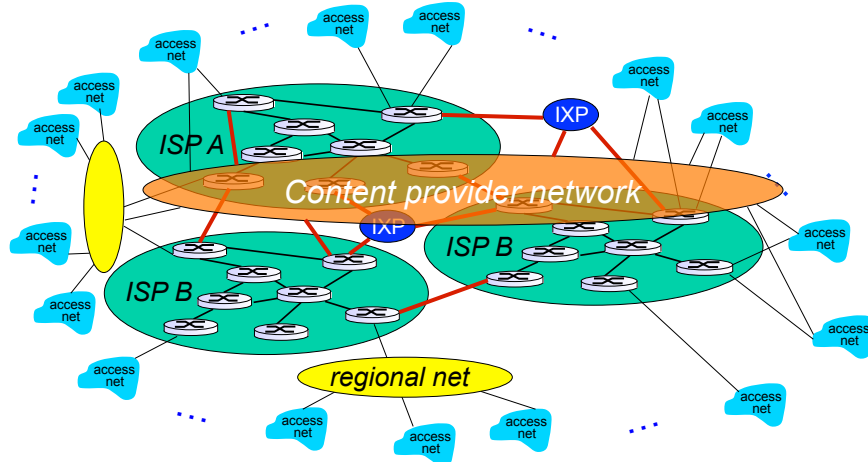
Internet I0I: a network of networks

... and regional networks may arise to connect access nets to ISPs



Internet I0I: a network of networks

... and content providers may run their own network, to bring services, content close to end users



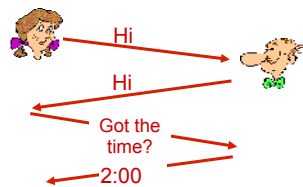
Internet I/O: protocols

human protocols:

- “what’s the time?”
- “I have a question”
- introductions

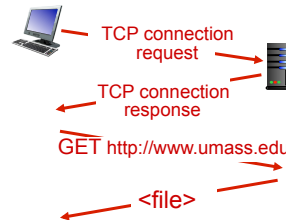
... specific messages sent

... specific actions taken when messages received



network protocols:

- hardware, software rather than humans
- define *format, meaning of messages* sent/received, *actions taken* on transmission/receipt
- all communication activity in Internet governed by protocols



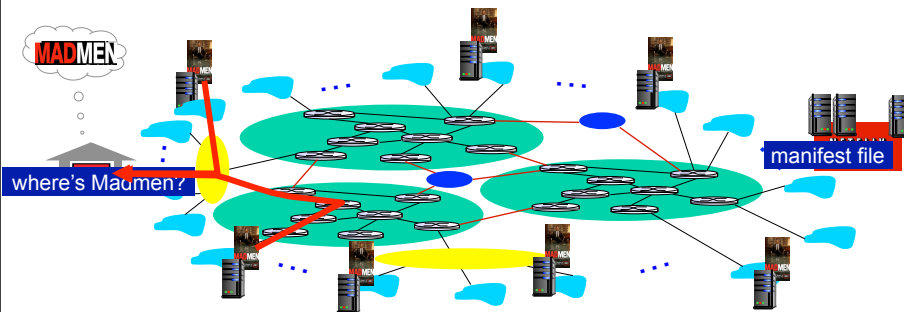
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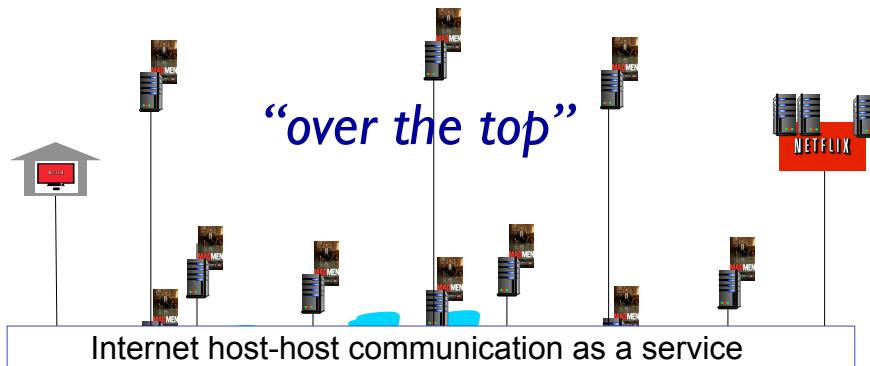
Challenge: can a network architecture designed for wired, host-host communication meet demands of today’s (and tomorrow’s) *content-oriented* mobile Internet?

Content Distribution Networks (CDNs)

- ❖ CDN: stores copies of content at CDN nodes
 - e.g. Netflix stores copies of MadMen
- ❖ subscriber requests content from CDN
 - directed to nearby copy, retrieves content
 - may choose different copy if network path congested



Content Distribution Networks (CDNs)

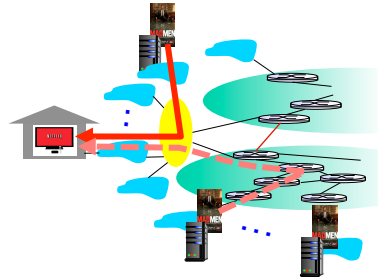


- OTT challenges:** coping with a congested Internet
- from which CDN node to retrieve content?
 - viewer behavior in presence of congestion?
 - what content to place in which CDN node?

CDNs: locating a “good” server

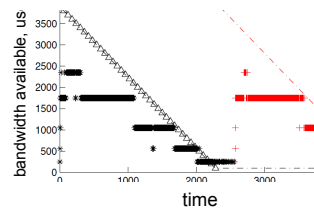
- ❖ CDNs (Akamai, Netflix): proprietary algorithms matching user with “nearby” server with high-bandwidth congestion-free path to client

- “maps” of Internet topology
- historical data (measurements) of Internet performance



Nygren, Sitaraman, Sun, "The Akamai network: a platform for high-performance internet applications," SIGOPS Oper. Syst. Rev. 44, 3 (August 2010)

- ❖ **client control:** request lower/higher quality video encoding, using less/more bandwidth
- **client control:** seek new server if performance drops below minimum



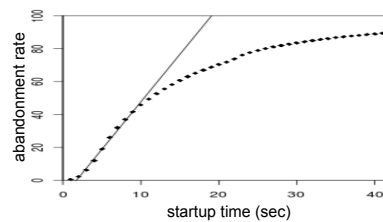
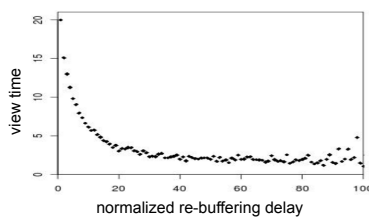
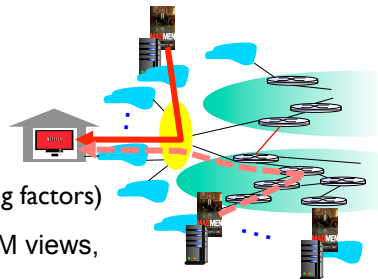
Adhikari, Guo, Hao, Varvello, Hilt, Steiner, Zhang, "Unreeling netflix: Understanding and improving multi-CDN movie delivery," IEEE INFOCOM, 2012.

CDN and user behavior

- ❖ **Question:** how does improved CDN performance affect user behavior?

- abandonment, engagement, repeat viewership
- quantitative, causal (avoid confounding factors)

- ❖ (big) data analysis: 6.7M users, 23M views, 100K unique videos

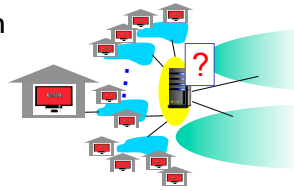


Video Stream Quality Impacts Viewer Behavior: Inferring Causality using Quasi-Experimental Designs, Shunmuga, Sitaraman, ACM Internet Measurement Conf. (IMC), Boston, MA, Nov 2012

CDN: content placement in servers

❖ *Question:* which content to place in which servers?

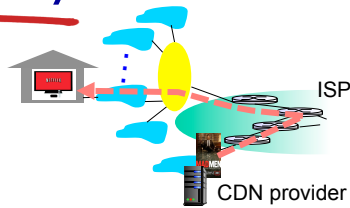
- Netflix: all servers identical
- predict future content access using ML techniques: content viewing, popularity ranking by neighbors?
 - populate server, make recommendation accordingly?
 - degree of difference among regional servers?



.... on going work (Dernbach, Kurose, Mahadevan) with Technicolor Research

DETOUR network neutrality

Question: to what extent can ISP regulate (or differentially charge) traffic flows within its network?



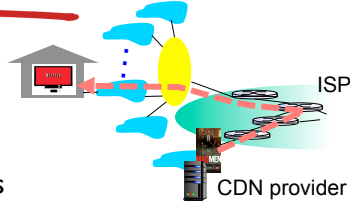
FCC “Open Internet Order” (2010):

- ❖ *[Transparency rule:]* ... shall publicly disclose accurate information ... sufficient for consumers to make informed choices ... and for content, application, service, and device providers to develop, market, and maintain Internet offerings
- ❖ *[No blocking rule:]* ... shall not block lawful content, applications, services, or non-harmful devices, subject to reasonable network management.
- ❖ *[Non discrimination rule:]* ... shall not unreasonably discriminate in transmitting lawful network traffic over a consumer’s broadband Internet access service.



network neutrality

- ❖ **2000-2005: deregulation** – ISPs re-defined as “information service” providers rather than “telecommunication service” providers
- ❖ **2005:** ISP (Madison River) FCC consent decree: stop dropping VoIP packets
- ❖ **2008:** FCC order against Comcast for throttling BitTorrent file sharing. Overturned 2010: FCC lacks sufficient statutorily-mandated responsibility under Title I
- ❖ **Feb. 2014:** Verizon v FCC: vacates *Open Internet Order*
- ❖ **March 2014:** Netflix, noting degrading performance to Comcast subscribers, agrees to “pay-to-peer” with Comcast



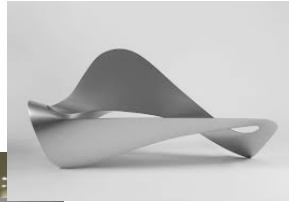
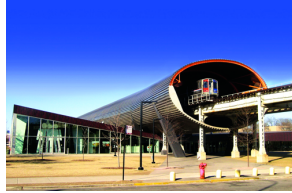
The future?

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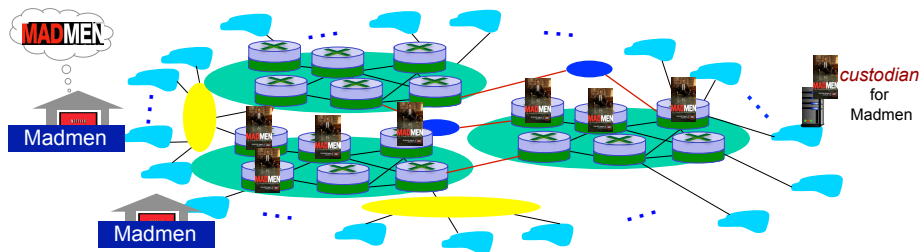
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Architecture: form follows function



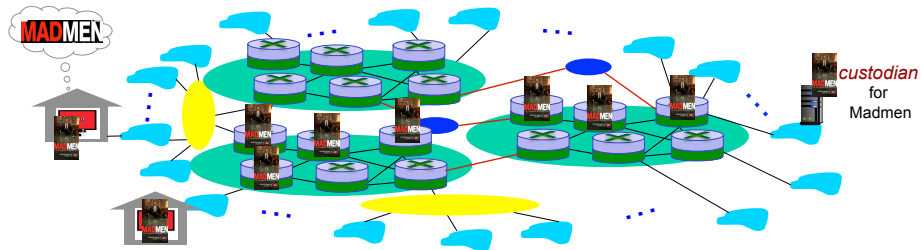
*so too in network
architecture?*

In-network caching



- ❖ routers equipped with massive storage
- ❖ client request routed towards custodian
- ❖ content returned, cached enroute
- ❖ client requests may be satisfied by in-network caches

In-network caching

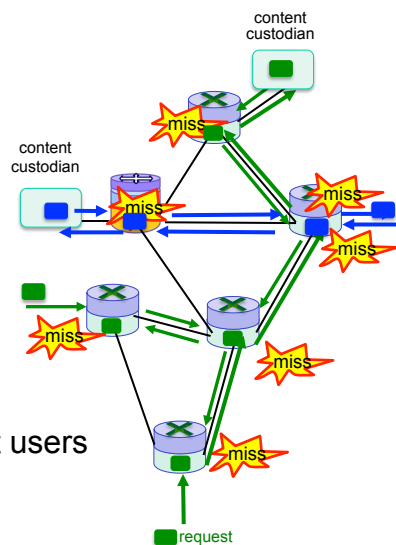


Many challenges!

- ❖ routing to custodian: based on name, or address?
 - how to search for content
- ❖ mobile users, content custodians
- ❖ cache management algorithms
- ❖ modeling, design, analysis of cache networks

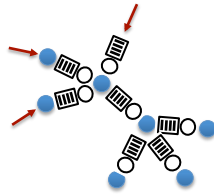
Cache networks

- ❖ client *requests content*
- ❖ request *routed* (e.g., shortest path) to known *content custodian*
- ❖ en-route to custodian, *caches* inspect request
 - *hit*: return local copy
 - *miss*: forward request towards custodian
- ❖ during content download, *store* in caches along path
- ❖ content requests from different users interact: *cache replacement*



Challenge: networks of caches

- ❖ *network effect: interaction* among content request/reply flows from different users:
 - ❖ *content replacement*: requested content by one user replaces content previously requested by others



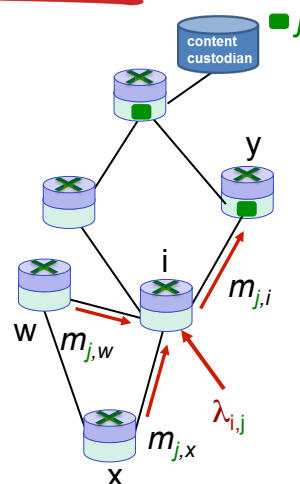
Packet-switching:
queueing networks
(Kleinrock, 1963)

Modeling a network of caches

- ❖ node i : *exogenous* (external) arrivals for content j : $\lambda_{i,j}$
- ❖ node i : *internal* arrivals (*miss stream*) for content j from downstream neighbors h : $m_{j,h}$
 - complex, correlated process
- ❖ $r_{i,j}$: aggregate rate of arrival requests at i for content j

$$r_{i,j} = \lambda_{i,j} + \sum_{\substack{\text{all downstream} \\ \text{neighbors, } h}} m_{j,h}$$

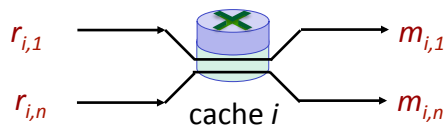
- ❖ ZDD: zero download delay assumption



Cache networks: approximate analysis

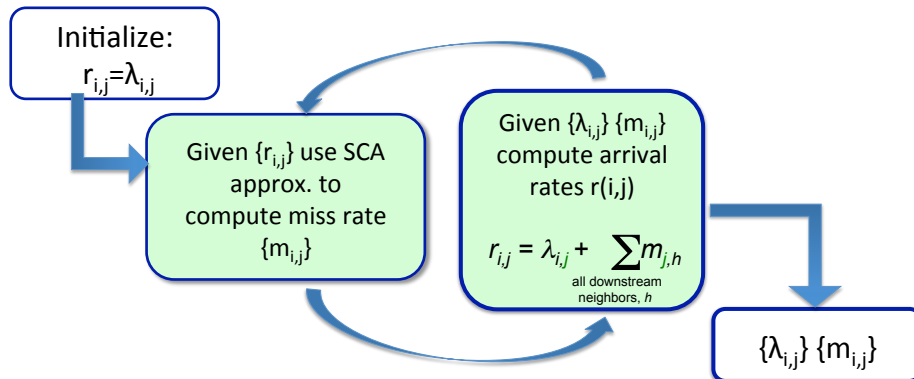
- ❖ **SCA**: standalone cache i approximation algorithm: given $r(i,j)$, compute miss rate for all content j
- ❖ Independence Reference Model (IRM) of incoming requests:

$$Pr(X_t = j \mid X_1, \dots, X_{t-1}) = Pr(X_t = j)$$
- ❖ SCA approximation algorithm for LRU: [Dan 1985]



But we need $\{(r_{i,j}, m_{i,j})\}$ for a *network* of caches

Fixed-point iteration



- ❖ fixed-point approximation relatively accurate: main source of error is independent reference model

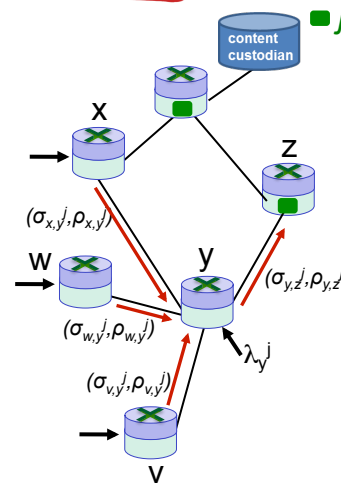
"Approximate Models for General Cache Networks," Elisha J. Rosensweig, Jim Kurose and Don Towsley, IEEE NFOCOM 2010

Cache networks: bounding analysis

Question: can we provably *bound* characteristics of request flows between caches?

- $\rho_{v,y}^j$: bound on average arrival rate of requests for content item j , from cache v to cache y
- $\sigma_{v,y}^j$: bound on burstiness of requests for content item j , from cache v to cache y

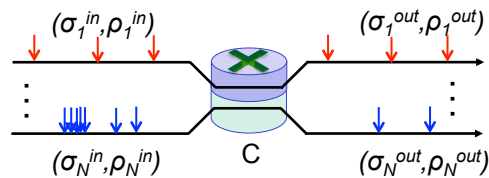
$$\int_{t_1}^{t_2} r_{y,z}^j(t) dt < (t_2 - t_1) \rho_{v,y}^j + \sigma_{v,y}^j$$



From $(\sigma_i^{in}, \rho_i^{in})$ to $(\sigma_i^{out}, \rho_i^{out})$

network calculus for cache flows:

- given input bounds, compute output bounds



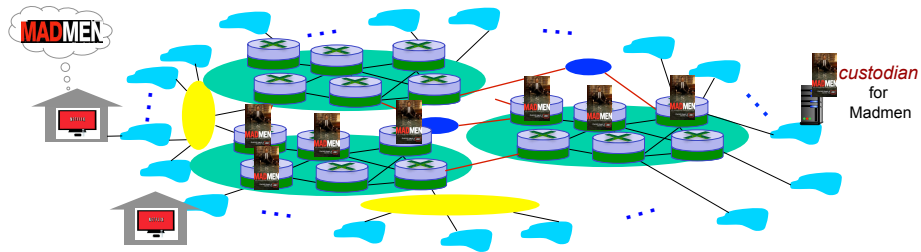
Theorem: For a cache of size C :

$$\rho_i^{out} = \min(\rho_i^{in}, M_i(\rho_1^{in}, \dots, \rho_n^{in}, C))$$

Can calculate ρ_i^{out} from $\{\rho_i^{in}\}$, σ_i^{out} also

E. Rosensweig, J. Kurose, "A Network Calculus for Cache Networks," IEEE Infocom 2013.
 E. Rosensweig, D. Menasche, J. Kurose, "On the Steady-State of Cache Networks," IEEE INFOCOM 2013
 (Best paper award)

In-network caching



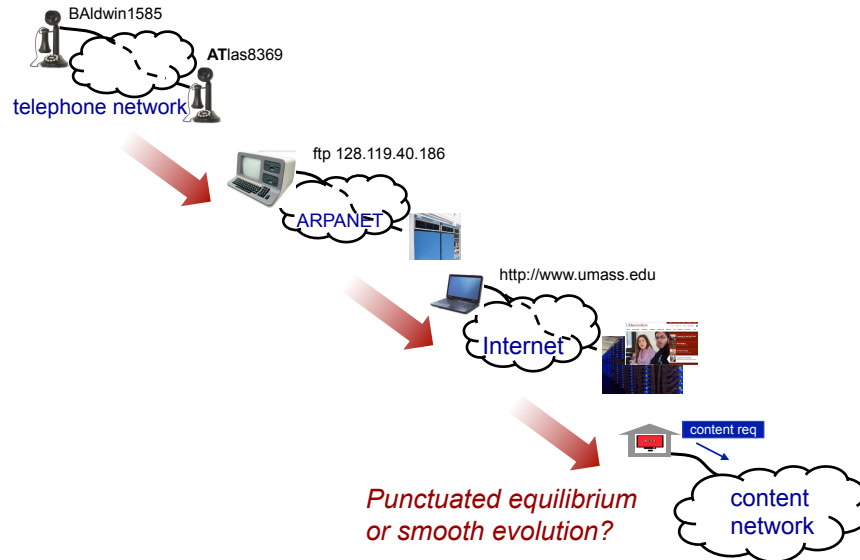
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Getting connected



Thanks !

*UMass
faculty
collaborators*

*... and staff
and friends*

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S. Mahadevan
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B. Woolf
M. Zink

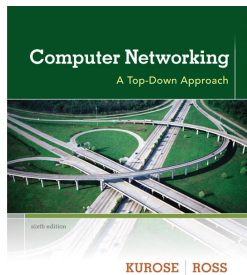


Thanks !

*UMass
students
and
postdocs*

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Y. Guo	E. Nahum	S. Singh	T. Zhu
D. Gyllstrom	G. Neglia	S. Singh	Z. Zhu

To learn more:



- ❖ CS 290nw: A networked world
- ❖ CS 453: Computer networks